# **Basic Elements of Programming Languages**

Ronghui Gu Spring 2020

Columbia University

<sup>\*</sup> Course website: https://www.cs.columbia.edu/ rgu/courses/4115/spring2019

<sup>\*\*</sup> These slides are borrowed from Prof. Edwards.

# What is a Programming Language?

A programming language is a notation that a person and a computer can both understand.

- It allows you to express what is the **task** to compute
- It allows a computer to execute the computation task

**Language Specifications** 

## How to Define a Language

When designing a language, it's a good idea to start by sketching forms that you want to appear in your language as well as forms you do not want to appear.

```
int avg(int a, int b)
{
  return (a + b) / 2;
}
```

Examples

```
a int vg(int a,
{
   return (a; + b)
{
```

Non-Examples

# How to Define a Language

- An official documents, with **informal** descriptions.
- An official documents, with **formal** descriptions.
- A reference implementation, e.g., a compiler.

Some language definitions are sanctioned by an official standards organization, e.g., C11 (ISO/IEC 9899:2011).

```
int compare()
{
   int a[10], b[10];
   if (a > b)
     return true;
   return false;
}
```

## **Aspects of Language Specifications**



- Syntax: how characters combine to form a program.
- **Semantics**: what the program *means*.
- **Pragmatics**: common programming idioms; programming environments; the standard library; ecosystems.

#### **Syntax**

#### Syntax is divided into:

- Microsyntax: specifies how the characters in the source code stream are grouped into tokens.
- Abstract syntax: specifies how the tokens are grouped into phrases, e.g., expressions, statements, etc.

#### Microsytax

Source program is just a sequence of characters.

```
int avg(int a, int b)
{
  return (a + b) / 2;
}
```

```
i n t SP a v g ( i n t SP a , SP i n t SP b ) NL { NL SP SP r e t u r n SP ( a SP + SP b ) SP / SP 2 ; NL } NL
```

## Microsytax

```
int avg(int a, int b)
{
  return (a + b) / 2;
}
```

Token	Lexemes	Pattern (as regular expressions)
ID	avg, a, b	letter followed by letters or digits
KEYWORD	int, return	letters
NUMBER	2	digits
OPERATOR	+, /	+, /
PUNCTUATION	;,(,),{,},	;,(,),{,},





## **Lexical Analysis Gives Tokens**

```
int avg(int a, int b)
{
  return (a + b) / 2;
}

int avg((int a , int b)) { (return (a + b)) / 2; }
```

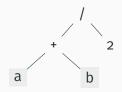
 Throw errors when failing to create tokens: malformed numbers (e.g., 23f465#g) or invalid characters (such as non-ASCII characters in C).

#### **Abstract Syntax**

Abstract Syntax can be defined using Context Free Grammar.

```
expr :
expr OPERATOR expr
( expr )
NUMBER
```

Expression (a + b)/2 can be parsed into an AST:

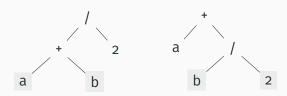


#### **Abstract Syntax**

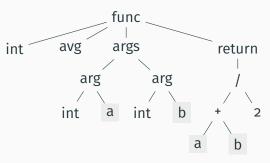
Abstract Syntax can be defined using Context Free Grammar.

```
expr :
    expr OPERATOR expr
| ( expr )
| NUMBER
```

Ambiguous! What about a + b/2?



## Syntax Analysis Gives an Abstract Syntax Tree



```
int avg(int a, int b)
{
   return (a + b) / 2;
}
```

 Syntax analysis will throw errors if "}" is missing. Lexical analysis will not.

#### **Semantics**

- **Static Semantics**: deals with legality rules—things you can check before running the code (compile time), e.g., type, scope, for some languages.
- **Dynamic Semantics**: deals with the execution behavior; things that can only be known at runtime, e.g., value.

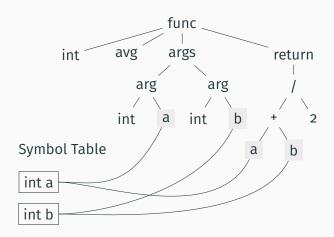
#### **Static Semantics**

We can use inference rules to define semantics, e.g., type:

$$\frac{\text{expr} : \textbf{int}}{\text{NUMBER} : \textbf{int}} \qquad \frac{\text{expr} : \textbf{int}}{(\text{expr}) : \textbf{int}}$$

 $\frac{\mathsf{expr}_1 \, : \, \mathbf{int} \quad \mathsf{expr}_2 \, : \, \mathbf{int}}{\mathsf{expr}_1 \, \, \mathsf{OPERATOR} \, \mathsf{expr}_2 \, : \, \mathbf{int}}$ 

# Semantic Analysis: Resolve Symbols; Verify Types



## **Dynamic Semantics**

We can use inference rules to define semantics, e.g., value:

$$\frac{\mathbf{eval}(\mathsf{expr}) = n}{\mathbf{eval}(\mathsf{expr}_1) = n \cdot \mathbf{eval}((\mathsf{expr}_2) = n}$$
 
$$\frac{\mathbf{eval}((\mathsf{expr}_1) = n_1 \quad \mathbf{eval}((\mathsf{expr}_2) = n_2 \quad (n_1 + n_2) = n}{\mathbf{eval}((\mathsf{expr}_1) + \mathsf{expr}_2) = n}$$

## **Dynamic Semantics**

#### Consider the integer range:

$$\frac{\mathsf{wrap}(\mathsf{NUMBER}) = n}{\mathsf{eval}(\mathsf{NUMBER}) = n} \qquad \frac{\mathsf{eval}(\mathsf{expr}) = n}{\mathsf{eval}((\mathsf{expr})) = n}$$

$$\frac{\mathbf{eval}(\mathsf{expr}_1) = n_1 \quad \mathbf{eval}(\mathsf{expr}_2) = n_2 \quad \mathbf{wrap}(n_1 + n_2) = n}{\mathbf{eval}(\mathsf{expr}_1 \ + \ \mathsf{expr}_2) = n}$$

**Programming Paradigms** 

## **Programming Paradigms**

A programming paradigm is a style, or "way," of programming. Some languages make it easy to write in some paradigms but not others.

## **Imperative Programming**

An imperative program specifies how a computation is to be done: a sequence of statements that update state.

```
result = []
    i = 0
    numStu = len(students)
start:
    if i >= numStu goto finished
    name = students[i]
    nameLength = len(name)
    if nameLength <= 5 goto nextOne
    addToList(result, name)
nextOne:
    i = i + 1
    goto start
finished:
    return result
```

## **Structured Programming**

A kind of imperative programming with clean, goto-free, nested control structures. Go To Statement Considered Harmful by Dijkstra.

```
result = []
for i in range(len(students)):
    name = students[i]
    if len(name) > 5:
        addToList(result, name)
print(result)
```

## **Procedural Programming**

Imperative programming with procedure calls.

```
def filterList (students):
    result = []
    for name in students:
        if len(name) > 5:
            addToList(result, name)
    return result

print(filterList(students))
```

# **Object-Oriented Programming**

An object-oriented program does its computation with interacting objects.

```
class Student:
  def init (self, name):
    self.name = name
    self.department = "CS"
def filterList (students):
    result = []
    for student in students:
        if student.name. len () > 5:
            result.append(student.name)
    return result
print(filterList(students))
```

# **Declarative Programming**

A declarative program specifies what computation is to be done. It expresses the logic of a computation without describing its control flow.

```
select name
from students
where length(name) > 5
```

# **Functional Programming**

A functional program treats computation as the evaluation of mathematical functions and avoids side effects.

```
def isNameLong (name):
    return len(name) > 5

print(
    list(
    filter(isNameLong, students)))
```

# **Functional Programming**

#### Using lambda calculus:

```
print(
  list(
  filter(lambda name: len(name), students)))
```

# **Functional Programming**

#### Using function composition:

```
\begin{array}{c} {\rm compose(\,print\,,\,\,list\,\,,\,\,\,filter\,*(lambda\,\,name:\,\,len\,(name)\,\,>\,\,5))}\\ {\rm (\,students\,)} \end{array}
```

<sup>\*</sup> A variant of the built-in filter.